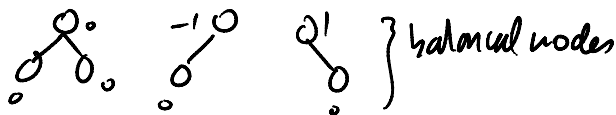


- AS64...
- Text's AVL tree works by recalculating the height of a node following insertion
- An alternate approach is to maintain balance info at each node & return the height change following insertion
- Changes to text's AVL implementation:
 - first notice the code can be streamlined somewhat by removal of the double-left & double-right routines (but calls them doubleWithLeftChild, doubleWithRightChild)
 - this will "compress" insertion code



~~template~~ <typename T>

~~void~~ Tree<T>: insert(const T& x, Node*& t)

{

int delta_h = 0;

if (t == NULL) { t = new Node(x, NULL, NULL); delta_h = 1; }

else if (x < t->data) {

~~insert(x, t->left);~~

~~if (height(t->left) - height(t->right) == 2) {~~

if (insert(x, t->left)) { // height of left subtree increased

t->bal--;

if (t->bal == -1) delta_h = 1;

else if (t->bal == -2) {

if (x < t->left->data)

rotate_left(t)

else {

~~double_left(t)~~

rotate_right(t->left)

rotate_left(t)

}

if (t->left->bal == 1)

~~if (x > t->left->data)~~

rotate_right(t->left)

rotate_left(t)

return

}

else if (x > t->data) {

insert(x, t->right)

if (height(t->right) - height(t->left) == 2)

```

    if (x < t → right → delta)
        rotate-left(t → right)
    rotate-right(t)
}

```

↖ rewrite

```

else if (x > t → delta) ?
    if (insert(x, t → right) ?
        t → bal++;
        if (t → bal == 1) delta-h = 1;
        else if (t → bal == 2) ?
            if (t → right → bal == -1)
                rotate-left(t → right)
            rotate-right(t)
        }
    }
}

```

```

}
else ; // no-op

```

~~t → height = max(height(t → left), height(t → right)) + 1~~

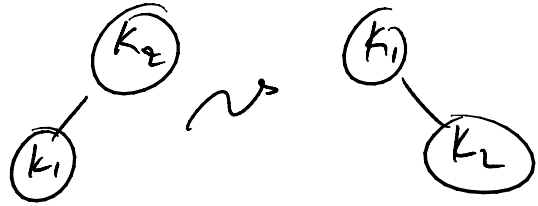
return delta-h;

}

before:

void Tree <T>:: rotate_left (Node * & k_L)

{
Node * k₁ = k_L → left;
k_L → left = k₁ → right;



k₁ → right = k_L;

k₂ → height = ...

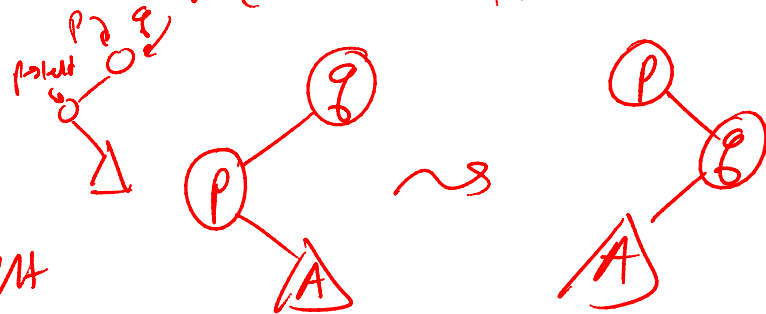
k₁ → height = ...

k_L = k₁;

}

void Tree <T>:: rotate_left (Node * & p)

{
Node * q = p → left;
p = p → left
q → left = p → right



p → right = q

q → height++;

if (p → height < 0) q → height -= p → height

p → height++;

if (q → height > 0) p → height += q → height

}

- A → ...

```
void Tree<T>::rotate - rot (Node & nP)  
{  
    // mirror of rotate - left  
}
```

Deletion:

Alg: if node is a leaf remove it

else
replace it with the minimum of
the right subtree then erase that node

After deletion, retrace path back
up the tree to the root, adjusting
the balance factors as needed

d.j.

like

insert($v, t \rightarrow left$)

```

t → bal -- ;
if (t → bal == 0) delta_h = 1;
else if (t → bal == -2) ?
    if (t → left → bal == 1)
        rotate_right(t → left)
    rotate_left(t)
if (t → bal == 0) delta_h = 1
}

```

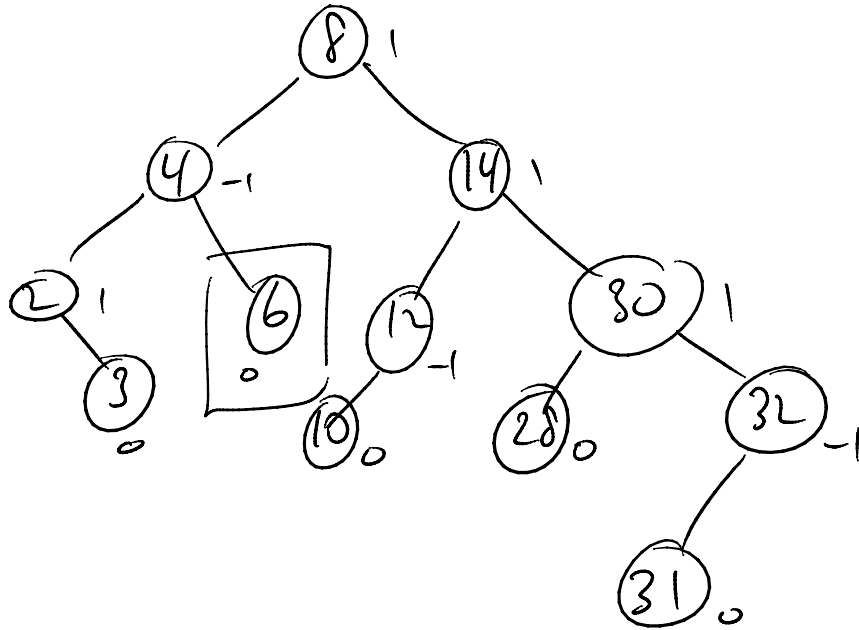
-the remainder of the delete routine precedes the
above in two tests to rework down to either
left or right subtree to perform deletion

recall the normal BST delete:

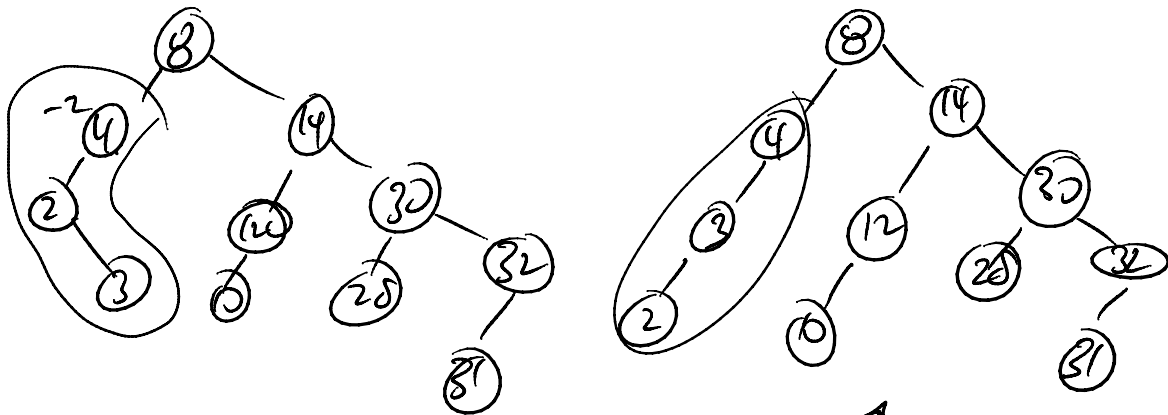
if ($x < t \rightarrow \text{data}$)
 erase ($x, t \rightarrow \text{left}$)
else if ($x > t \rightarrow \text{data}$)
 erase ($x, t \rightarrow \text{right}$)

- these now have to be expanded in manner
analogous to above deletion of interior node

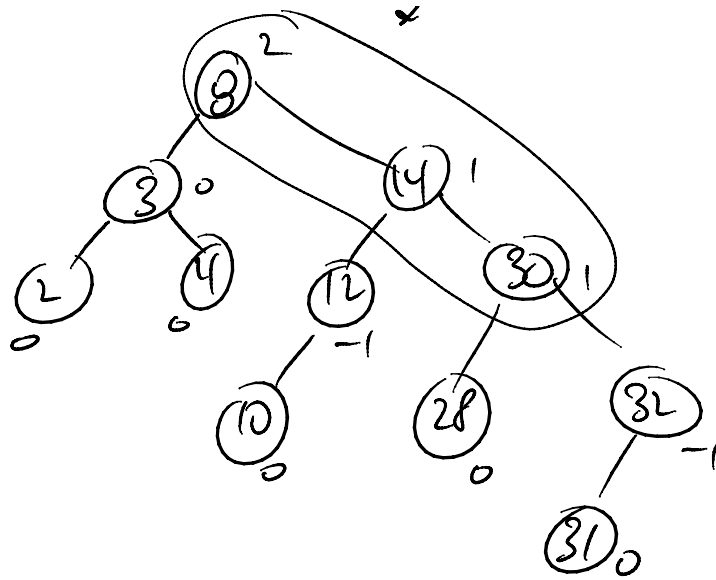
ex. input: insert 6, 4, 2, 8, 10, 12, 14, 32, 30, 28, 3, 31 ;



delete 6 :



rotate - right (t → left)
rotate - left (t)



rotate_right(8)

