

## Administrativa

1. My office hours: M 3-4 pm
2. Refrain from sending me code  
if you need help with C++, i.e.  
compiling, get a tutor

Last time: 2 methods for solving recurrence relations

Method 1: telescoping sum, ex. for mergesort

$$T(1) = 1$$

$$T(n) = 2T\left(\frac{n}{2}\right) + \underbrace{n}_N$$

divide by  $n$

$$\frac{T(n)}{n} = \frac{2T\left(\frac{n}{2}\right)}{n} + 1 = \frac{2}{n} \cdot T\left(\frac{n}{2}\right) + 1 = \frac{T\left(\frac{n}{2}\right)}{\frac{n}{2}} + 1$$

telescope  
out  
trying to  
reduce  
 $n$  down to 1

$$\frac{T\left(\frac{n}{2}\right)}{\frac{n}{2}} = \frac{T\left(\frac{\frac{n}{2}}{2}\right)}{\frac{\frac{n}{2}}{2}} + 1 = \frac{T\left(\frac{n}{4}\right)}{\frac{n}{4}} + 1$$

$$\frac{T\left(\frac{n}{4}\right)}{\frac{n}{4}} = \frac{T\left(\frac{\frac{n}{4}}{2}\right)}{\frac{\frac{n}{4}}{2}} + 1$$

⋮

⋮

$$n=2 \quad \frac{T(\sim)}{2} = \frac{T(1)}{1} + 1$$

add up

add up

add it all up (telescopic sum)

$$\frac{T(n)}{n} + \cancel{\frac{T(n/2)}{n/2}} + \cancel{\frac{T(n/4)}{n/4}} + \dots + \frac{T(2)}{2} =$$

$$\cancel{\frac{T(n/2)}{n/2}} + 1 + \cancel{\frac{T(n/4)}{n/4}} + 1 + \cancel{\frac{T(n/8)}{n/8}} + 1 + \dots + \frac{T(1)}{1} + 1$$

$$\boxed{\frac{T(n)}{n}} = \sum 1 + \frac{T(1)}{1} = \lg n + \frac{T(1)}{1}$$

why  $\lg n$ ?

$$T(n) = n \lg n + n \in O(n \lg n)$$

let's say  $n=8$

$$\frac{T(n)}{n} = \frac{T(8)}{8} \quad 2^3$$

$$\frac{T(n/2)}{2} = \frac{T(4)}{4} \quad 2^2$$

$$\frac{T(n/4)}{4} = \frac{T(2)}{2} \quad 2^1$$

$$\frac{T(n/8)}{8} = \frac{T(1)}{1} \quad 2^0$$

Method 2: recursive back substitution

$$T(1) = 0$$

$$T(n) = 2 + \left(\frac{n}{2}\right) + n - 1$$

$$\text{level 1: } T(n) = n - 1 + 2T\left(\frac{n}{2}\right)$$

$$\begin{aligned} \text{level 2: } &= n - 1 + 2\left(\frac{n}{2} - 1 + 2T\left(\frac{n}{4}\right)\right) \\ &= n - 1 + n - 2 + 4T\left(\frac{n}{4}\right) \end{aligned}$$

$$\begin{aligned} \text{level 3: } &= n - 1 + n - 2 + 4\left(\frac{n}{4} - 1 + 2T\left(\frac{n}{8}\right)\right) \\ &= n - 1 + n - 2 + n - 4 + 8T\left(\frac{n}{8}\right) \end{aligned}$$

$$\begin{aligned} \text{level 4: } &= n - 1 + n - 2 + n - 4 + 8\left(\frac{n}{8} - 1 + 2T\left(\frac{n}{16}\right)\right) \\ &= n - 1 + n - 2 + n - 4 + n - 8 + 16T\left(\frac{n}{16}\right) \end{aligned}$$

$$\text{level } k: = \sum_{i=0}^{k-1} n - \sum_{i=0}^{k-1} 2^i + 2^k T\left(\frac{n}{2^k}\right)$$

look up sheet with handy formulas

$$\sum_{i=0}^n k = k(n+1)$$

$$\sum_{i=0}^n 2^i = 2^{n+1} - 1$$

want this to be  $T(1)$

$$n(k-1+1) - (2^{k-1+1} - 1) + 2^k T\left(\frac{n}{2^k}\right)$$

$$= kn - 2^k + 1 + 2^k T\left(\frac{n}{2^k}\right)$$

$$\text{let } 2^k = n \Rightarrow \underline{k = \lg n}$$

$$\begin{aligned}
 &= n \lg n - n + 1 + n \left( \frac{n}{a} \right) \\
 &= n \lg n - n + 1 + \cancel{n} \cdot 1 \\
 &= n \lg n - n + 1 = O(n \lg n - n) \in O(n \lg n)
 \end{aligned}$$

# Handy formulae

$$\sum_{k=0}^n k = \frac{n(n+1)}{2}$$

$$\begin{array}{c} 1+2+3+\dots+n \\ n+(n-1)+(n-2)+\dots+1 \\ \hline (n+1)+(n+1)+\dots+(n+1) \end{array}$$
$$\sum_{i=0}^n i = \frac{n(n+1)}{2}$$

$$\sum_{i=0}^n 2^i = 2^{n+1} - 1$$

$$\sum_{i=0}^n i 2^i = (n-1)2^{n+1} + 2$$

$$\sum_{i=0}^{\lg n - 1} 2^i = n - 1$$

$$\sum_{i=0}^{\lg n - 1} n \lg \frac{n}{2^i} = \frac{n \lg^2 n}{2} + \frac{n \lg n}{2}$$

For homework:

review Quicksort (worst case  $\neq$  best case)

$$T(0) = T(1) = 1$$

$$T(n) = T(n-1) + cn$$

$$T(n) = cn + 2T\left(\frac{n}{2}\right)$$

$$k=0 \quad = cn + T(n-1)$$

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