

Secrets of the kd-Tree revealed

① The function `object`
(aka the functor)

~~Point Axis Compare~~
`point_c`

originally based on
text's version on p.34-35
("Comparator")

class point {
public:
} → an obj with just one member

a) bool operator()
(const point_t & p1,
const point_t & p2)
const

{ return (p1[axis] < p2[axis])

can be 0, 1, 2

↳ private;
int axis;

}; // point_c class def.

where is this set/initialized?

↓
in point_c constructor:

c) point_c (int in axis = 0) :
axis (in axis) {} ;

d) one more version
of operator () N
needed for point_t &
pointer

bool operator ()

(const point_t & const & p1,
const point_t & const & p2)

const

{ return ... }

with $point_c$ defined,

$kdtree_t < point_t,$

$point_p,$

$point_c > kdtree;$

② KdTree insertion & min & max

KdTree_t (T, P, C) ::

insert (Knode_t * &t,

std::vector<P> &x,

const T &min,

const T &max,

int d)

}

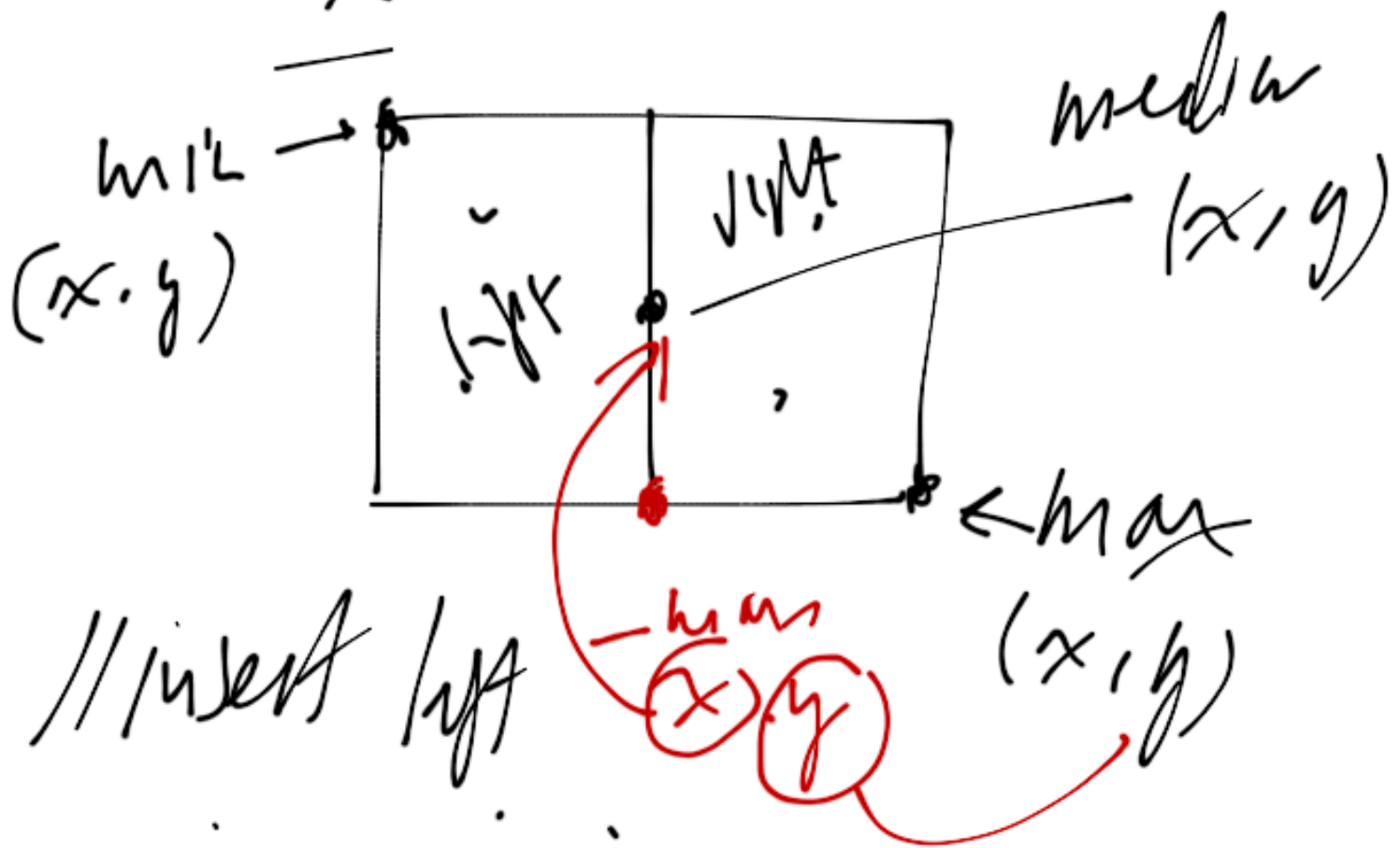
// local VWS (at center
tree level)

T — min, — max;

// What got sent
to next tree level
in recursive call to
insert

// code to find
median, split &
ans (0, 1, 2)

e.g. median is midpoint
 in vertical axis, sorted
 on axis



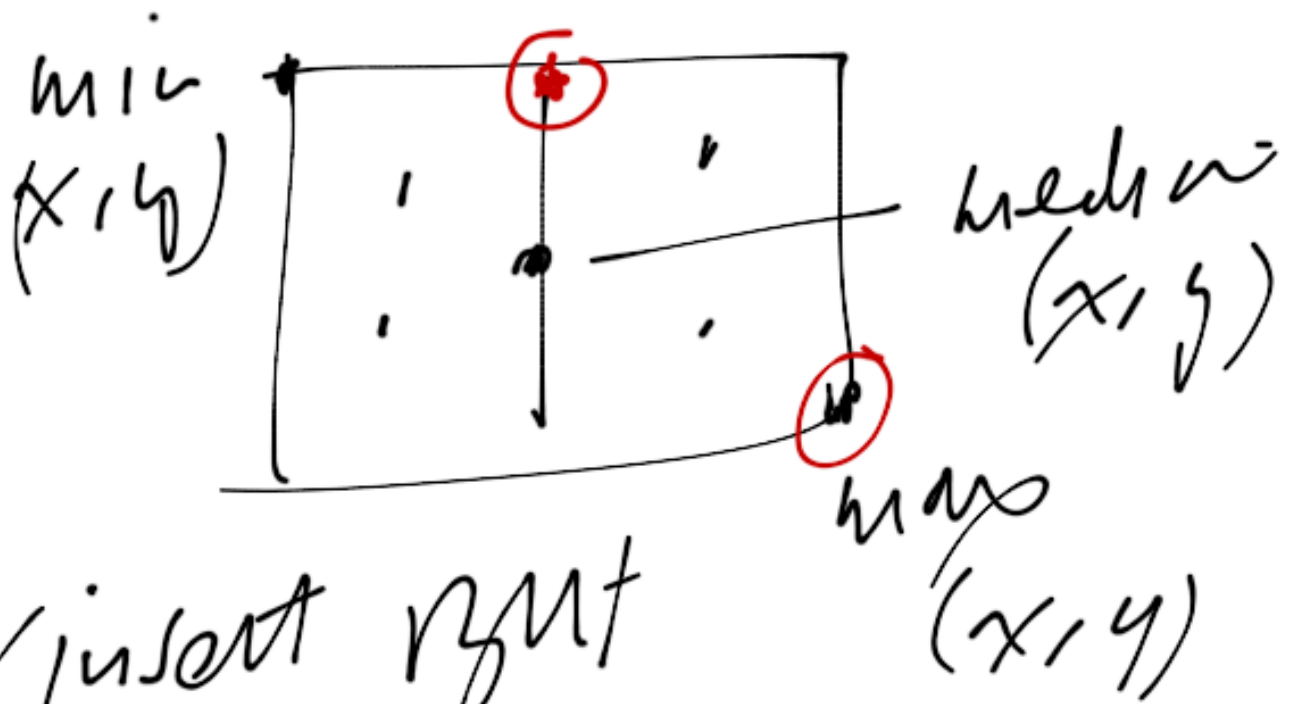
// insert left

~~min~~
~~(x, y)~~

— min = min;

— max = max;

— max[axis] = (x, median[axis])



// insert point

- min = min

- max = max

- $\min(\text{axis}) = (\text{xmedian})(\text{axis})$

③ NN query (log n)

$dist = g \cdot distance(t \rightarrow data)$

↑
query
point t

↑
tree node
ptr to
point t

if ($dist < r$) {

$r = dist;$

$p = t \rightarrow data;$

} // traverse down L, R sub
tree

KNN query: (logic)

$\text{dist} = g.\text{distance}(t \rightarrow \text{data})$

~~$\hat{y}(\text{dist} < r)$~~

~~$r = \text{dist}$~~

~~$P = t \rightarrow \text{data}$~~

~~$\therefore P$ is now a vector!~~

P_0	P_1	P_2	\dots	P_k
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for k closest points

— need to find k^{th} - closest

distance: just look at P_k

assuming list is sorted

p_k is just $p.back()$

std::vector

if (dist < q.distance(q.back()))

r = dist

put $t \rightarrow$ data

on list, maintaining

list of size K

} CAVEATS: list may
have fewer than K , must be sorted

$dist = g \cdot distance(t \rightarrow data)$
// do we keep this t or not?

if size of list $< K$ {
// insert t into
sorted list

if (list is empty ||

$dist > g \cdot distance(p.back())$)

push $t \rightarrow data$ at end

else

insert $t \rightarrow data$
into proper position

}

else if list has K or if
already

done only in else (once we know K)

!only insert t onto
list if dist < last
node's on list

if (dist < g.distance(p.back()))

insert t → data
into proper position

if (p.size() > K)

p.pop_back()

r = g.distance(p.back())