

- kd-tree insertion
- kd-tree NA, Knn queries
- basic insertion mechanism (p. 551)

```

public:
    void insert ( typename P
                 (point_t*)
                 std::vector<P> & X,
                 const T & min,
                 const T & max )
    { root = insert (root, X, min, max, 0);
  }
  
```

private:

Kdnode\_t \*root;

Kdnode\_t \*insert(

Kdnode\_t \*Q,

std::vector<P> Q,

const T &, const T &,

int);

in kdtree.cpp:

```
template < typename T,  
          typename P,  
          typename C >
```

```
typename Kdtree_t < T, P, C > ::  
    Kdnode_t *
```

```
Kdtree_t < T, P, C > :: insert (  
    Kdnode_t * &t,  
    std::vector < P > &x,  
    const T &min, const T &max,  
    int d)
```

{ int axis =  
x.empty() ? 0 :  
d % x[0] → dim();

in photon +  
class, handed  
to 3

int m = 0; // median  
index

P median;  
T - min, - max;

// sampling vol. of subspace

std::vector<P> left, right;

typename std::vector<P>::  
iterator iE;

if (x.empty()) return NULL;

// find median by sorting  
// (not very efficient)

sort(x.begin(), x.end(),

cc(axis)); *functor*  
→ #include <algorithm> *in photo, h*

// get median

$m = x.size() / 2;$

// create left & right subarray

for (int i = 0; i < (int)  
x.size();  
i++)

if (i < m)

left.push\_back(  
x[i])

else if (i > m)

right.push\_back(x[i])  
the median = x[m];

// create new node

Kdnode\_t \* node =

new Kdnode\_t (median,  
min,  
max,  
NULL,  
NULL,  
axis)

recursively add left tree

min =

max =

min[axis] =

node → left = insert(node,  
left, min,  
max, d+1)

recursively add right tree

min =

max =

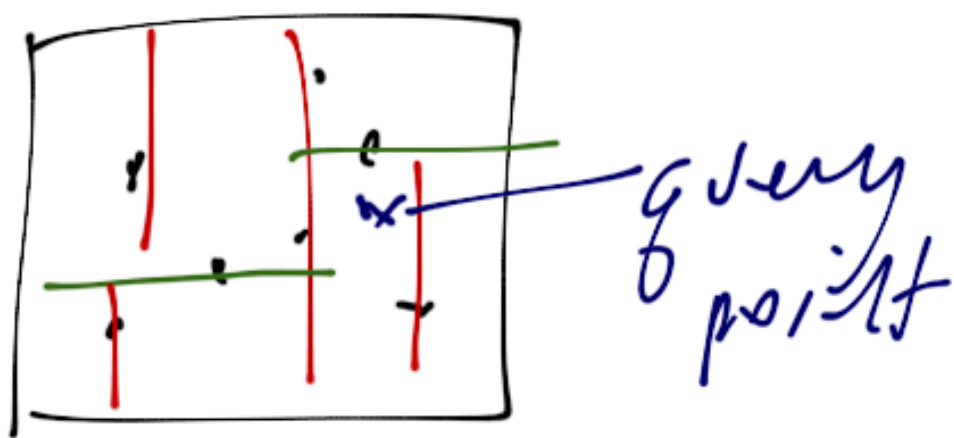
min[axis] =

node → right = insert(node,  
right, min,  
max, d+1);

return node;



- Nearest-neighbor query



- the query itself:

$input: q$ , the query node  
(fake photon),

$t$ , the node

(root to start with),

$Q, R$ , dist. thus far to nearest

(dist set to  $\infty$  initially)

\*  $u$  node, or  $PQ$

(pointer reference to  
nearest node)



if (dist < r) {

r = dist;

p = t → data;

} (as we descend tree  
test against each node  
we encounter

// traverse down "closed"  
side of tree

axis = t → axis;

if (g[axis] <= (∀ t → data)  
[axis])

}  
nn(t → left, g, p, r);

// as we return, check to  
see if circle def. by g & r  
intersect farther side of tree  
→

if ( $g[axis] + r >$   
 $(t \rightarrow data)[axis]$ )  
     $h_u(t \rightarrow right, g, p, r)$

} else {

$h_u(t \rightarrow right, g, p, r)$

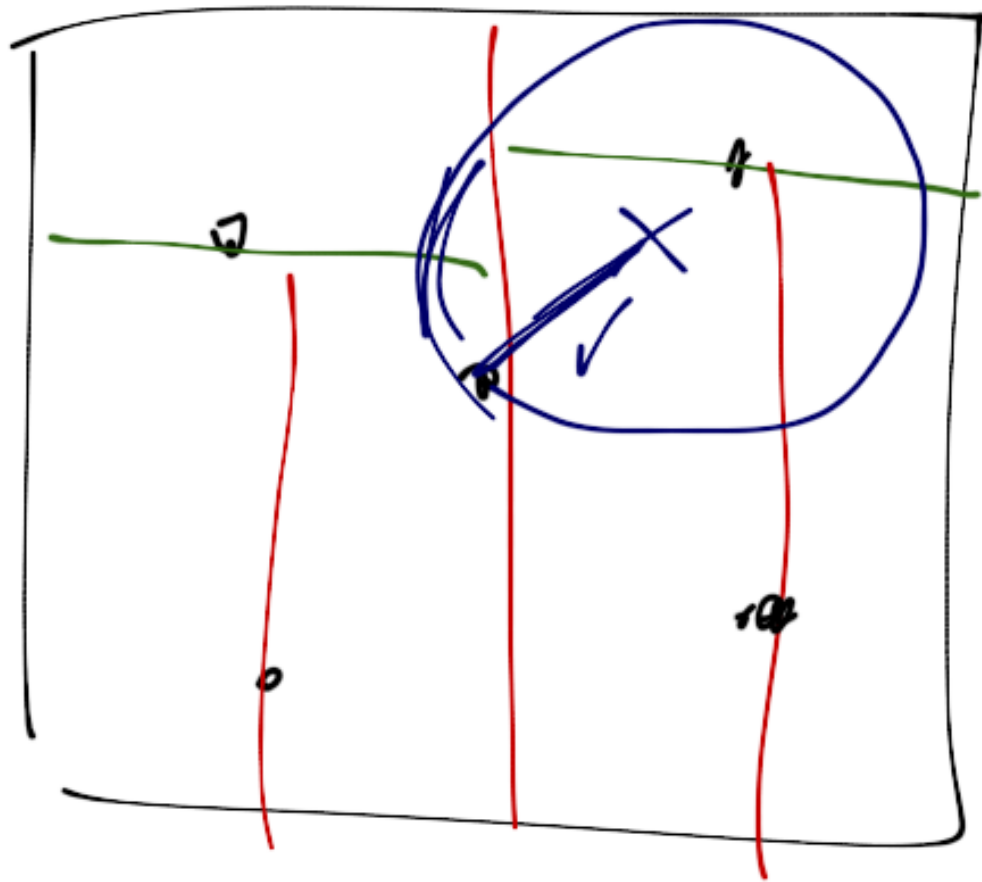
if ( $g[axis] - r <=$

$(t \rightarrow data)[axis]$ )

$h_u(t \rightarrow left, g, p, r)$

} }

- see Think Nguyen's  
lecture



-  $k$ -nearest neighbors:

- instead of only  
searching within a  
circle whose radius

is "closest distance yet",

search within a circle

whose radius is  $k^{\text{th}}$  closest  
yet found. UNTIL  $k$  POINTS



HAVE BEEN FOUND,

KEEP DISTANCE AT  $\infty$

- otherwise similar to  
4n query, just keep  
a sorted list of  $k$   
points found thus far

∴  
dist = g . distance (t → data)

what we find: a vector of photo pts

if (int) p.size() < k

if (p.empty())

(dist > g . distance (p.back()))

p.push\_back(t → data)

else {

// insert into list

```
for (pit = p.begin();
```

```
    pit != p.end();
```

```
    pit++) {
```

```
    p.insert (pit, l,
```

```
              t → data)
```

// insert node into sorted

list.

---

$V = g.distance(p.begin())$