

Quicksort

(with) indices

Alg (pseudocode, sort of):
if (right > left)
quicksort(array, left, right) {

// pick a pivot

pivot = partition(array,
left, right,
pivot)

also an index, value
(can change)

quicksort(array, left,
pivot - 1)

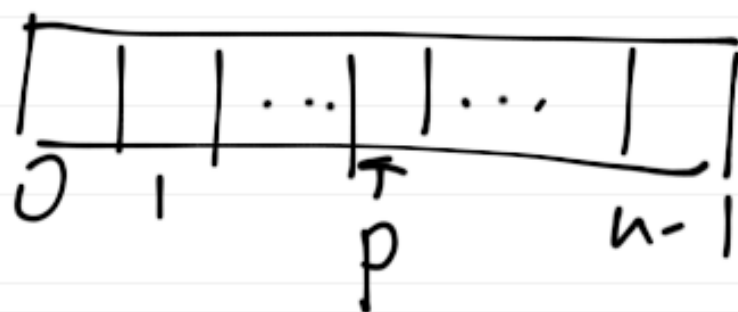
quicksort(array, pivot + 1,
right)

}

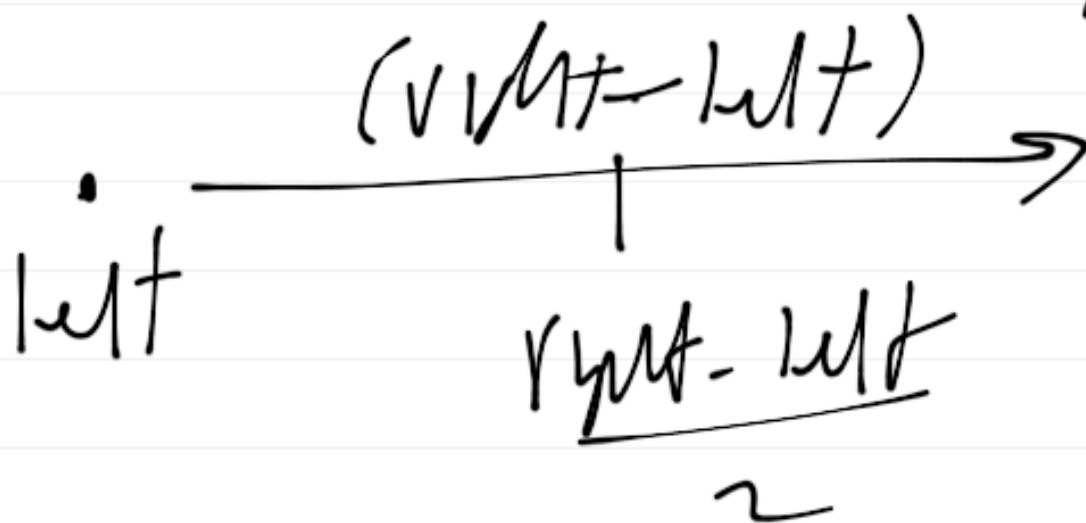
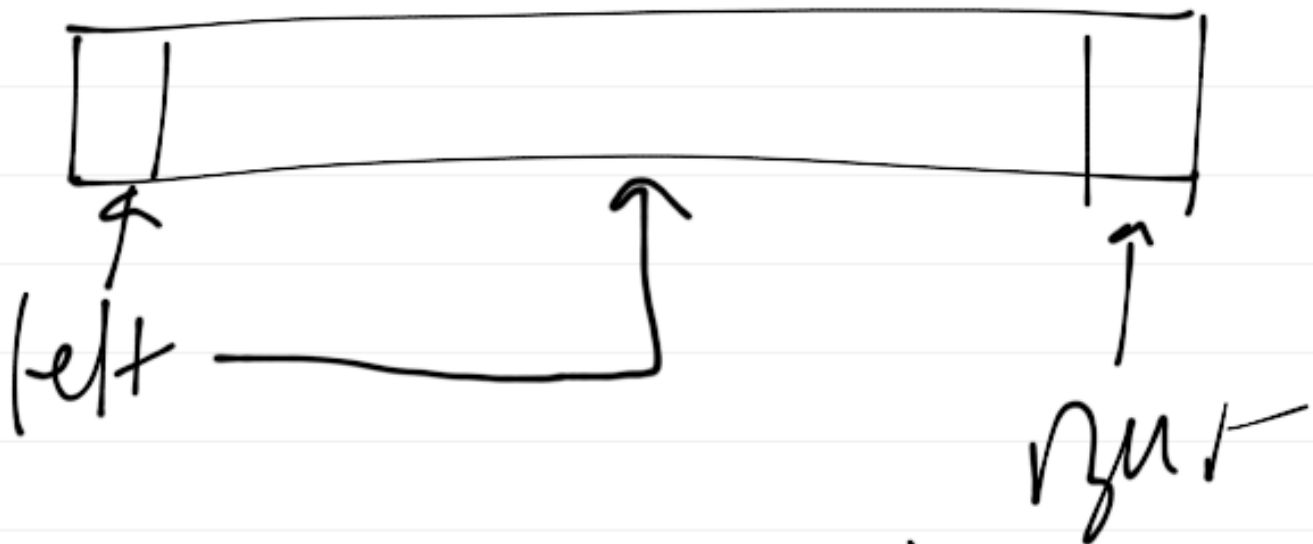
- picking a pivot:

ideally, pivot should be

index of median element



$$\text{pivot} = \text{left} + (\text{right} - \text{left}) / 2$$



Analysis:

$$T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) +$$

$$T(\text{partition})$$

partition (array, left, right, pivot)

1. store data at pivot location

the_pivot = array[pivot]

2. move pivot element to end
of array

swap(array[pivot], array[right])

3. go thru array, from left to right

moving values \leq the_pivot

to the left of pivot index,
advancing left index

```
for (i = left; i < right; i++) {  
    if (array[i] <= the_pivot) {  
        swap(arr[left], arr[i])  
        left++;  
    }  
}
```

4. move pivot element to its
final location

```
swap(array[left], array[right])
```

5. return pivot's new
location

return (left)

(note: 'left' argument
must be pass-by-value)

3 7 8 5 2 1 9 5 4

left
i

pivot

right



3 7 8 4 2 1 9 5 5

i...

pivot



3 4 8 7 2 1 9 5 5



3 4 2 7 8 1 9 5 5



3 4 2 1 8 7 9 5 5



3 4 2 1 5 7 8 9 5

at end of for loop:

3 4 2 1 5 7 8 9 5
left

last step puts pivot in left location

3 4 2 1 5 5 9 8 7

left list

right list

Note: Quicksort won't
degenerate to bubblesort
provided pivot is chosen
as the median element
is $O(n)$ time

↓
linear time median
select

Quickest worst case analysis

$$T(n) = T(i) + T(n-i-1)$$

$$+ \underbrace{cn}$$

moving elements
into position

$i = |S_1|$ the no. of elements
in one side of the

array

worst case: $i = 0$

Solve for $T(n)$

$$k=0$$

$$T(n) = cn + T(n-1)$$

$$k=1$$

$$= cn + (c(n-1) + T(n-2))$$

$$= cn + cn - c + T(n-2)$$

$$k=2$$

$$= cn + cn - c + (c(n-2) + T(n-3))$$

$$= (n + n + n - c - 2c + T(n-3))$$

$$k=3$$

$$= (n + cn + cn - c - 2c$$

$$+ (c(n-3) + T(n-4)))$$

$$= 4cn - c \sum_{k=0}^n k + T(n-4)$$

$$k=0$$

$$= c(k+1)n - \left(c \sum_{k=0}^n k + T(n-k+1) \right)$$

$$\sum_{k=0}^n k = \frac{n(n+1)}{2}$$

$$T(n) = c(k+1)n - \frac{cn(n+1)}{2}$$

$$+ T(n-k+1)$$

let $k=n$

$$T(n) = c(n^2+n) - \frac{cn^2}{2}$$

$$- \frac{cn}{2} + T(1)$$

$$\frac{cn^2}{2} - \frac{cn}{2} \in O(n^2)$$